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Discrete Convex Analysis III: Algorithms for Discrete Convex Functions

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Contents of Part III

Algorithms for Discrete Convex Functions

- A1.** Minimization (General)
- A2.** M-convex Minimization
- A3.** L-convex Minimization
- A4.** M-convex Intersection
- A5.** Submodular Maximization

A1.

Minimization (General)

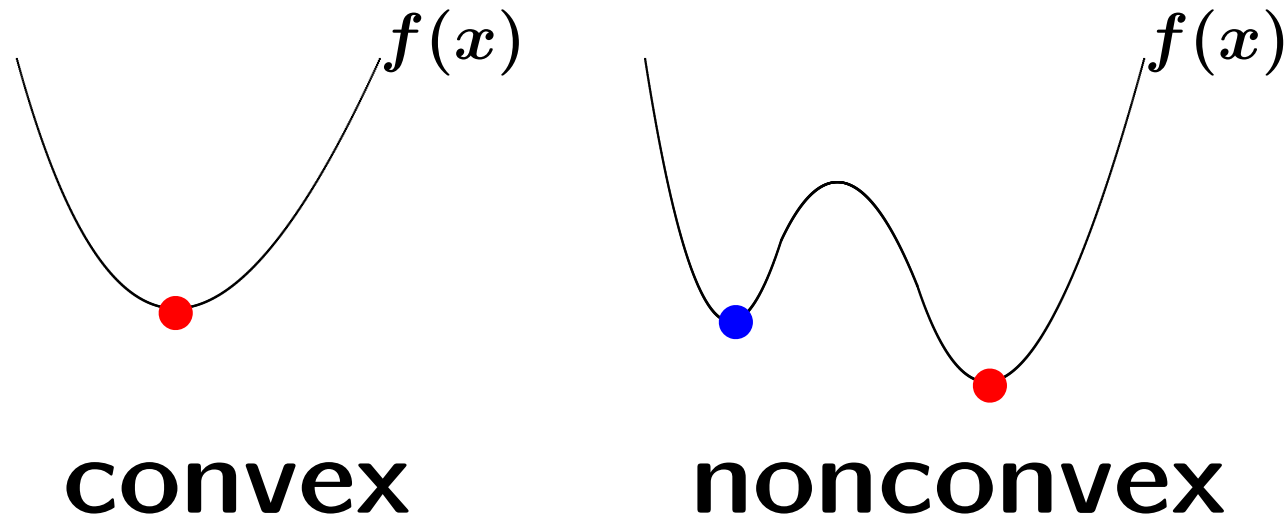
Optimality Criterion

Descent Method

Scaling and Proximity

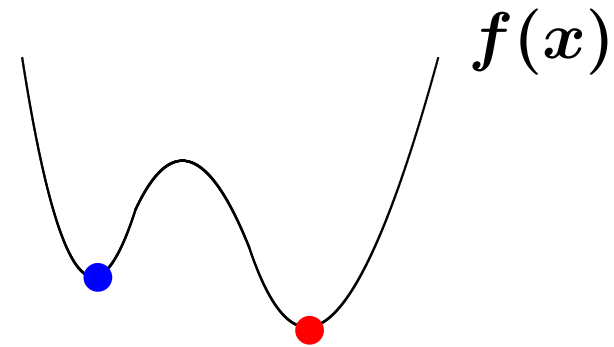
Optimality Criterion

Global opt vs **Local opt**



Local opt wrt neighborhood

Descent Method



S0: Initial sol x^*

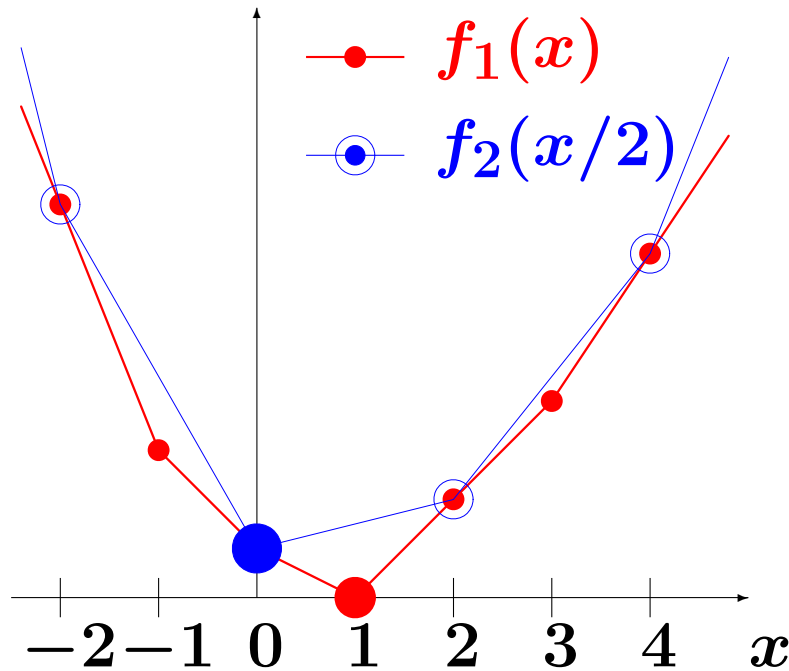
S1: Minimize $f(x)$ in **nbhd** of x^* to obtain x^\bullet

S2: If $f(x^*) \leq f(x^\bullet)$, return x^* (**local opt**)

S3: Update $x^* := x^\bullet$; go to S1

.....What is **nbhd** ?

Scaling and Proximity



Proximity theorem:

True minimum ● exists

in a neighborhood of

a scaled local minimum ●

⇒ efficient algorithm

Facts in DCA:

- Scaling preserves L-convexity
- Scaling does NOT preserve M-convexity
- Proximity thms known for L-conv and M-conv

A2.

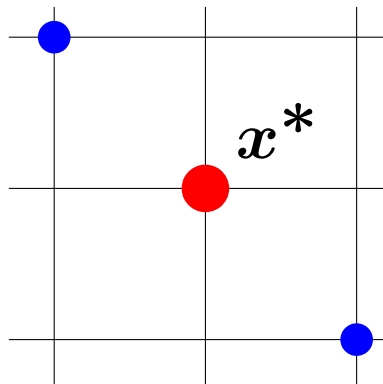
M-convex Minimization

Local vs Global Opt (M-conv)

Thm : $f : \mathbb{Z}^n \rightarrow \mathbb{R}$ M-convex (Murota 96)

x^* : global opt

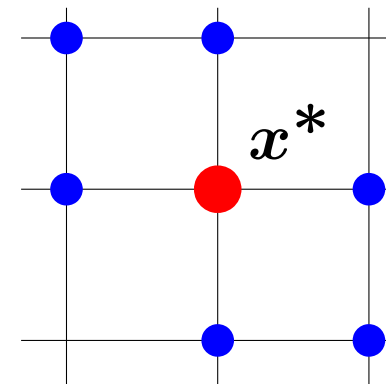
\iff local opt $f(x^*) \leq f(x^* - e_i + e_j) \quad (\forall i, j)$



Ex: $x^* + (0, 1, 0, 0, -1, 0, 0, 0)$

Can check with n^2 fn evals

For M ∇ -convex fn \Rightarrow



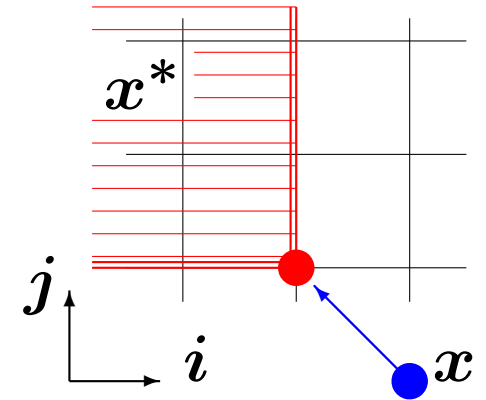
Steepest Descent for M-convex Fn

S0: Find a vector $x \in \text{dom } f$

S1: Find (i, j) that minimizes $f(x - e_i + e_j)$

S2: If $f(x) \leq f(x - e_i + e_j)$, stop
(x : minimizer)

S3: Set $x := x - e_i + e_j$
and go to S1



Minimizer Cut Thm

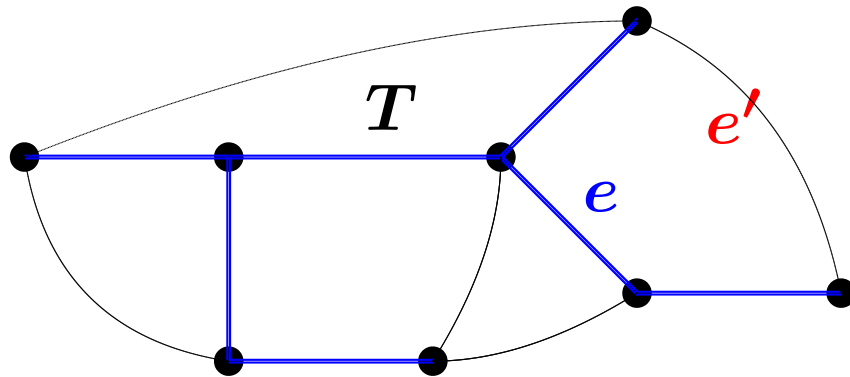
(Shioura 98)

\exists minimizer x^* with $x^*(i) \leq x(i) - 1$, $x^*(j) \geq x(j) + 1$

\Rightarrow Murota 03, Shioura 98, 03, Tamura 05

- Dress–Wenzel’s alg for valuated matroid
- Kalaba’s alg for min spanning tree

Min Spanning Tree Problem



length $d : E \rightarrow \mathbb{R}$

total length of T

$$\tilde{d}(T) = \sum_{e \in T} d(e)$$

Thm

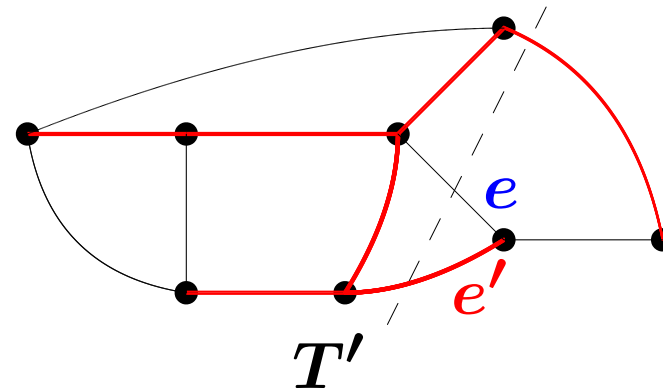
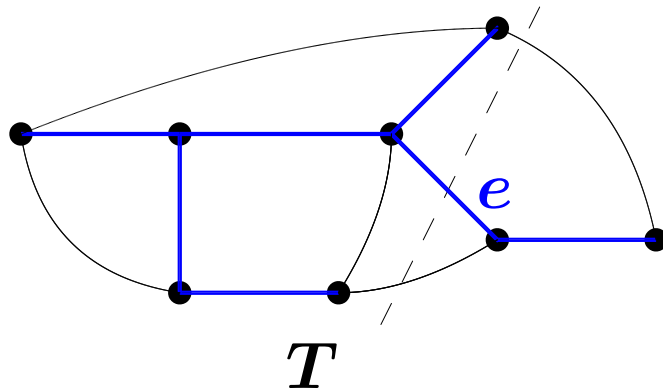
$$\begin{aligned} T: \text{MST} &\iff \tilde{d}(T) \leq \tilde{d}(T - e + e') \\ &\iff d(e) \leq d(e') \quad \text{if } T - e + e' \text{ is tree} \end{aligned}$$

Algorithm Kruskal's, Kalaba's

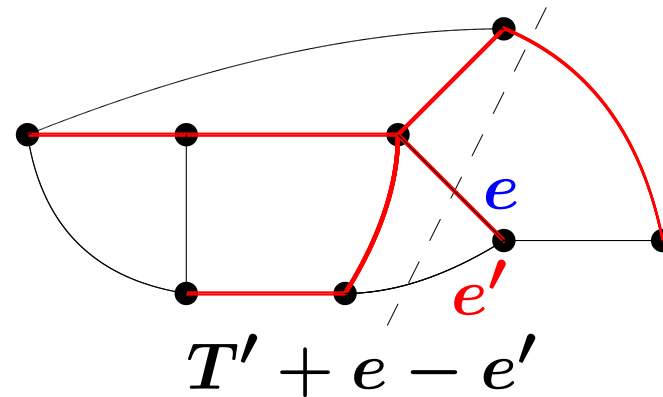
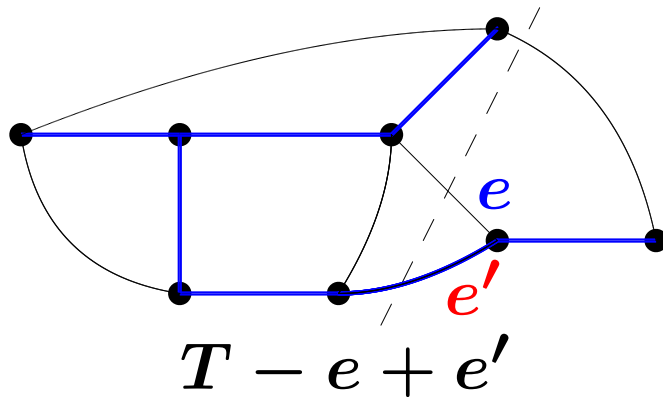
DCA view

- linear optimization on an M-convex set
- M-optimality: $f(x^*) \leq f(x^* - e_i + e_j)$

Tree: Exchange Property



Given pair
of trees



New pair
of trees

Exchange property: For any $T, T' \in \mathcal{T}$, $e \in T \setminus T'$
there exists $e' \in T' \setminus T$ s.t. $T - e + e' \in \mathcal{T}$, $T' + e - e' \in \mathcal{T}$

A3.

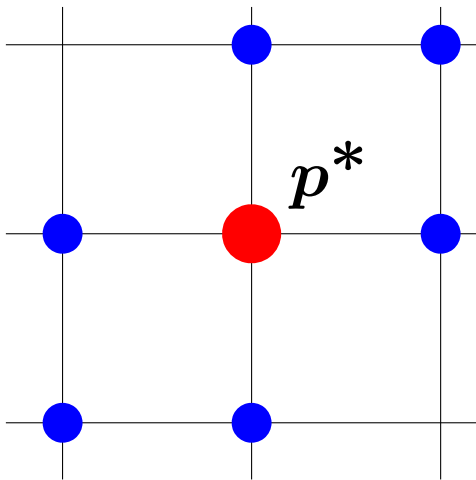
L-convex Minimization

Local vs Global Opt (L_{\square} -conv)

Thm : $g : \mathbb{Z}^n \rightarrow \mathbb{R}$ L_{\square} -convex (Murota 98,03)

p^* : global opt

\iff local opt $g(p^*) \leq g(p^* \pm q)$ ($\forall q \in \{0, 1\}^n$)



Ex: $p^* + (0, 1, 0, 1, 1, 1, 0, 0)$

$\iff \rho_{\pm}(X) = g(p^* \pm \chi_X) - g(p^*)$

takes min at $X = \emptyset$

Can check with n^5 (or less) fn evals
using submodular fn min algorithm
(Iwata-Fleischer-Fujishige, Schrijver, Orlin,
Lee-Sidford-Wong)

Steepest Descent for L^{\sharp} -convex Fn

(Murota 00, 03, Kolmogorov-Shioura 09, Murota-Shioura 14)

S0: Find a vector $p^{\circ} \in \text{dom } g$ and set $p := p^{\circ}$

S1: Find $\varepsilon = \pm 1$ and X that minimize $g(p + \varepsilon\chi_X)$

S2: If $g(p) \leq g(p + \varepsilon\chi_X)$, stop (p : minimizer)

S3: Set $p := p + \varepsilon\chi_X$ and go to S1

Thm:

(Murota-Shioura 14)

Termination exactly in $\mu(p^{\circ}) + 1$ iterations, where

$$\mu(p^{\circ}) = \min\{\|p^* - p^{\circ}\|_{\infty}^+ + \|p^* - p^{\circ}\|_{\infty}^- \mid p^* \in \arg \min g\}$$

$$\|q\|_{\infty}^+ = \max_i \max(0, q(i)), \quad \|q\|_{\infty}^- = \max_i \max(0, -q(i))$$

Monotone Steepest Descent for L^1 -convex Fn

S0: Find a vector $p^\circ \in \text{dom } g$ s.t

$\{q \mid q \geq p^\circ\} \cap \text{argmin } g \neq \emptyset$ and set $p := p^\circ$

S1: Find X that minimizes $g(p + \chi_X)$

S2: If $g(p) \leq g(p + \chi_X)$, stop (p : minimizer)

S3: Set $p := p + \chi_X$ and go to S1

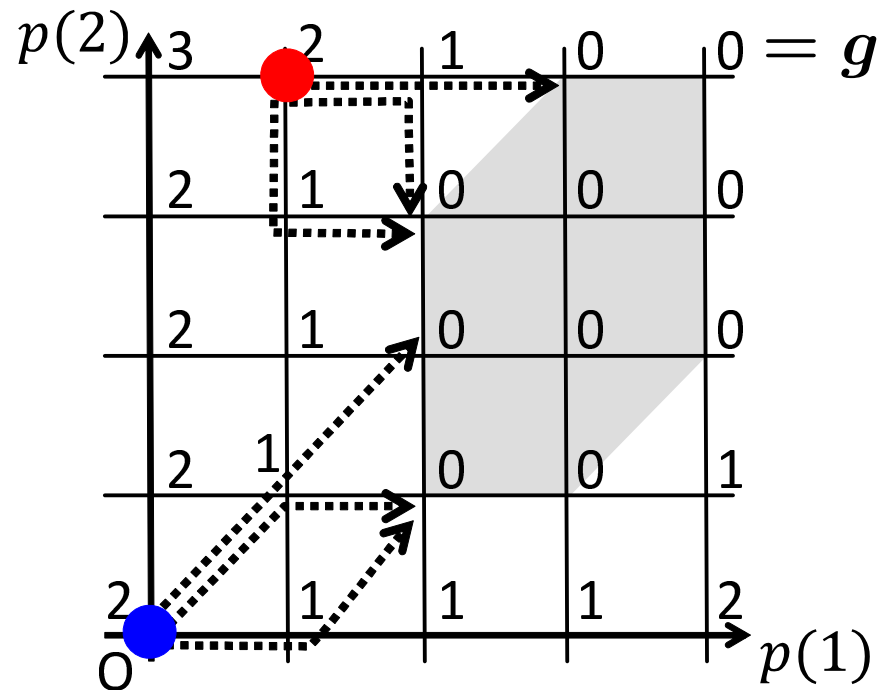
Thm: (Murota-Shioura 14)

Termination exactly in $\hat{\mu}(p^\circ) + 1$ iterations, where

$$\hat{\mu}(p^\circ) = \min\{\|p^* - p^\circ\|_\infty \mid p^* \in \text{argmin } g, p^* \geq p^\circ\}$$

\Rightarrow Application to ascending auction

Steepest Descent Path for L_1 -convex Fn



$$\mu(\mathbf{p}^\circ) = \hat{\mu}(\mathbf{p}^\circ) = 2$$

$$\|\mathbf{p}^\circ, \operatorname{argmin} g\|_\infty = 1$$

$$\mu(\mathbf{p}^\circ) = \hat{\mu}(\mathbf{p}^\circ) = 2$$

$$\|\mathbf{p}^\circ, \operatorname{argmin} g\|_\infty = 2$$

Shortest Path Problem (one-to-all)

one vertex (s) to all vertices, length $\ell \geq 0$, integer

Dual LP

$$\begin{aligned} & \text{Maximize } \sum p(v) \\ & \text{subject to } p(v) - p(u) \leq \ell(u, v) \quad \forall (u, v) \\ & \qquad \qquad p(s) = 0 \end{aligned}$$

Algorithm Dijkstra's

DCA view

- linear optimization on an L^{\natural} -convex set (in polyhedral description)
- Dijkstra's algorithm (M.-Shioura 12)
= steepest ascent for L^{\natural} -concave maximization
with uniform linear objective $(1, 1, \dots, 1)$

Optimality & Proximity Theorems

Func Class	Optimality	Proximity
L-convex	$f(x^*) \leq f(x^* + \chi_S) \quad (\forall S)$ $f(x^* + 1) = f(x^*)$ (M. 01)	$\ x^* - x^\alpha\ \leq (n-1)(\alpha-1)$ (Iwata-Shigeno 03)
M-convex	$f(x^*) \leq f(x^* - \chi_u + \chi_v)$ $(\forall u, v \in V)$ (M. 96)	$\ x^* - x^\alpha\ \leq (n-1)(\alpha-1)$ (Moriguchi-M.-Shioura 02)
L2-convex (L*L convol)	$f(x^*) \leq f(x^* + \chi_S) \quad (\forall S)$ $f(x^* + 1) = f(x^*)$	$\ x^* - x^\alpha\ \leq 2(n-1)(\alpha-1)$ (M.-Tamura 04)
M2-convex (M+M)	$f(x^*) \leq f(x^* - \chi_U + \chi_W)$ $(\forall U, W)$ (M. 01)	$\ x^* - x^\alpha\ \leq \frac{n^2}{2}(\alpha-1)$ (M.-Tamura 04)
integrally convex	$f(x^*) \leq f(x^* - \chi_U + \chi_W)$ $(\forall U, W)$ (Favati-Tardella 90)	<p style="text-align: center;">???</p>

$$\| \cdot \| = \| \cdot \|_\infty$$

A4.

M-convex Intersection

(Fenchel Duality)

Intersection Problem $(f_1 + f_2)$

Recall: $L^{\natural} + L^{\natural} \Rightarrow L^{\natural}, \quad M^{\natural} + M^{\natural} \not\Rightarrow M^{\natural}$

M-convex Intersection Algorithm:

Minimizes $f_1 + f_2$ for M^{\natural} -convex f_1, f_2

\Leftrightarrow Maximizes $f_1 + f_2$ for M^{\natural} -concave f_1, f_2
(submodular function maximization)

\Leftrightarrow Fenchel duality (min = max)

\Rightarrow Valuated matroid intersection (Murota 96)

\Rightarrow Weighted matroid intersection

(Edmonds, Lawler, Iri-Tomizawa 76, Frank 81)

M-convex Intersection: Min $[M^\natural + M^\natural]$

$M^\natural + M^\natural$ is NOT M^\natural

$f_1, f_2 : M^\natural$ -convex $(\mathbb{Z}^n \rightarrow \mathbb{R})$, $x^* \in \text{dom } f_1 \cap \text{dom } f_2$

(1) x^* minimizes $f_1 + f_2$ (Murota 96)

$\iff \exists p$ (certificate of optimality)

• x^* minimizes $f_1(x) - \langle p, x \rangle$ (M-opt thm)

• x^* minimizes $f_2(x) + \langle p, x \rangle$ (M-opt thm)

(2) $\text{argmin}(f_1 + f_2) = \text{argmin}(f_1 - p) \cap \text{argmin}(f_2 + p)$

(3) f_1, f_2 are integer-valued \Rightarrow integral p

M-concave Intersection: Max $[M^\natural + M^\natural]$

[Concave version]

$M^\natural + M^\natural$ is NOT M^\natural

$f_1, f_2 : M^\natural$ -concave $(\mathbb{Z}^n \rightarrow \mathbb{R})$, $x^* \in \text{dom } f_1 \cap \text{dom } f_2$

$f_1 + f_2$ is submodular

(1) x^* maximizes $f_1 + f_2$ (Murota 96)

$\iff \exists p$ (certificate of optimality)

• x^* maximizes $f_1(x) - \langle p, x \rangle$ (M-opt thm)

• x^* maximizes $f_2(x) + \langle p, x \rangle$ (M-opt thm)

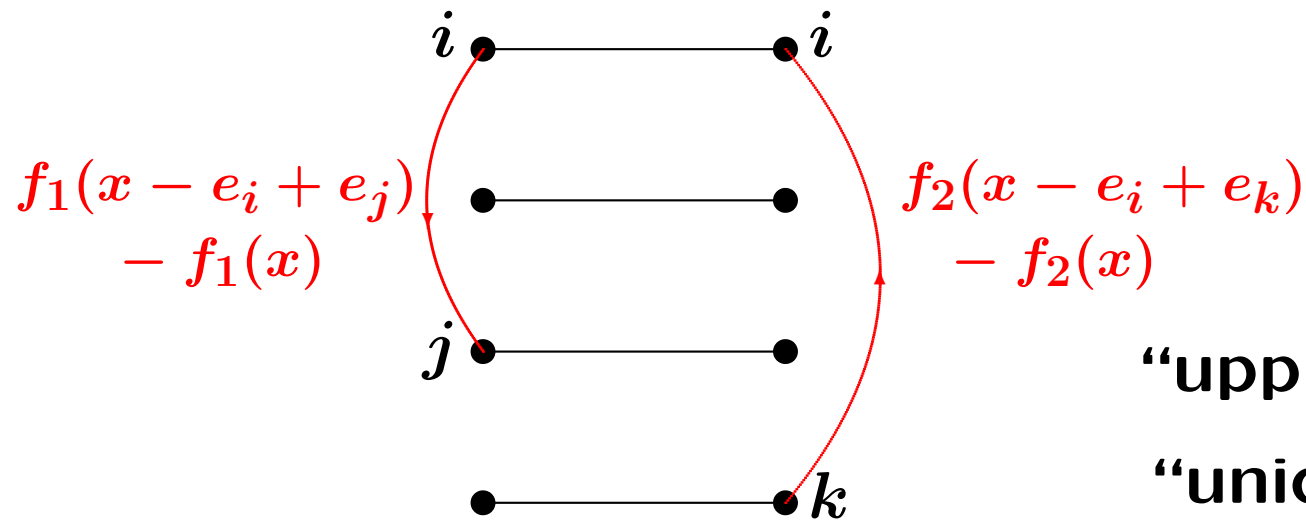
(2) $\text{argmax}(f_1 + f_2) = \text{argmax}(f_1 - p) \cap \text{argmax}(f_2 + p)$

(3) f_1, f_2 are integer-valued \Rightarrow integral p

M-convex Intersection Algorithms

Natural extensions of
weighted (poly)matroid intersection algorithms

Exchange arcs are weighted



“upper-bound lemma”

“unique-max lemma”

- cycle-canceling (Murota 96, 99)
- successive shortest path (Murota-Tamura 03)
- scaling (Iwata-Shigeno 03, Iwata-Moriguchi-M. 05)

Convolution

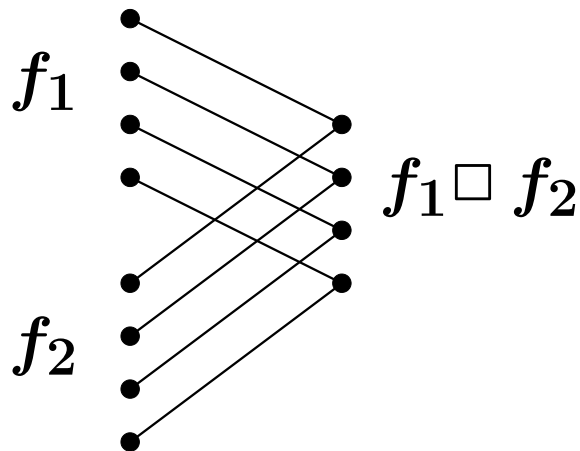
Convolutions of M^{\natural} -convex functions:

$$(f_1 \square f_2)(x) = \min_y (f_1(y) + f_2(x - y))$$

$$(f_1 \square f_2 \square f_3)(x), \quad (f_1 \square f_2 \square \cdots \square f_k)(x)$$

can be computed by **M-convex intersection algorithms**

cf. aggregated utility function

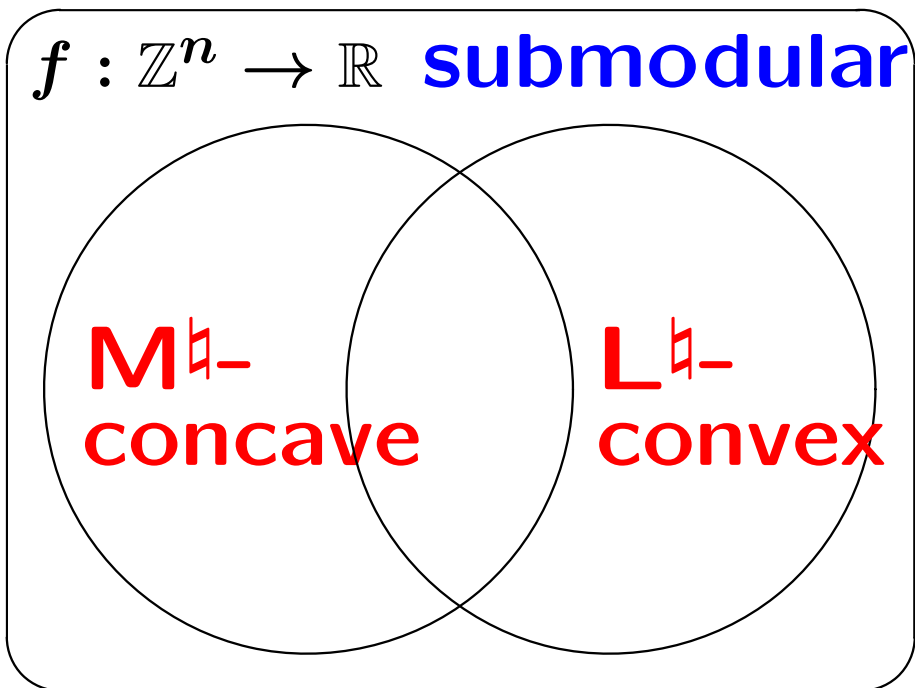


A5.

Submodular Maximization

Submodularity & Convexity in DCA

- M^\natural -concave function is submodular
- L^\natural -convex function is submodular



$$\text{sbm} + \text{sbm} \Rightarrow \text{sbm}$$

$$L^\natural + L^\natural \Rightarrow L^\natural$$

$$M^\natural + M^\natural \not\Rightarrow M^\natural$$

- Sum of M^\natural -concave fns is submodular

(i) M^\natural

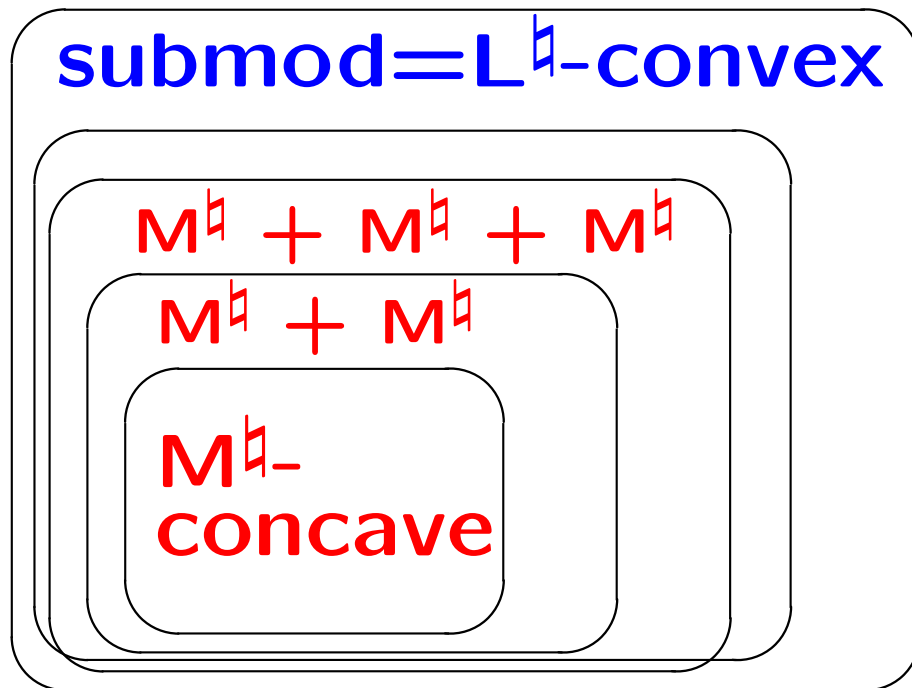
(ii) $M^\natural + M^\natural$

(iii) $M^\natural + M^\natural + M^\natural$

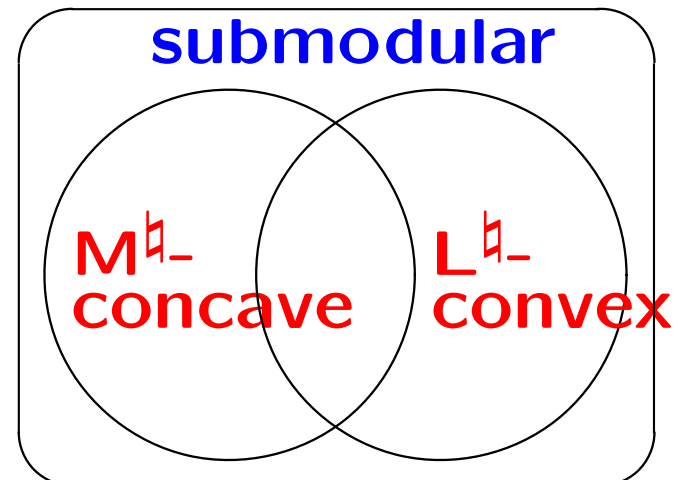
Submodular Set Function in DCA

- **Submodular** set func = **L[♠]-convex** on $\{0, 1\}^n$
- (Sums of) **M[♠]-concave** form nice subclasses

$$f : \{0, 1\}^n \rightarrow \mathbb{R}$$



$$f : \mathbb{Z}^n \rightarrow \mathbb{R}$$

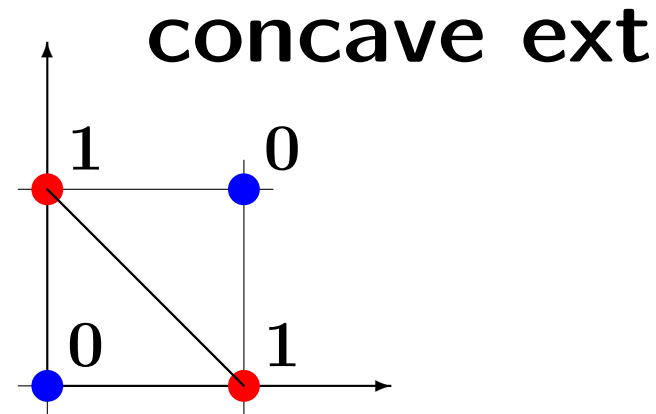
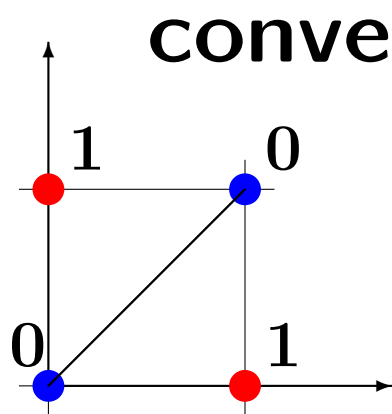


Set Function and Extensions

Set function \iff Function on $\{0, 1\}^n$

$$\rho(X) = \hat{\rho}(\chi_X)$$

Every set function $\rho : \{0, 1\}^n \rightarrow \mathbb{R}$ can be extended to convex/concave function



cf. Lovász extension

M[♯]-concave Set Functions

M[♯]-concave is **submodular** (NOT conversely)

M[♯]-concave forms a nice subclass for maximization

- $\mu(X) = \varphi(|X|)$ (φ : concave)
- $\mu(X) = \sum_{A \in \mathcal{T}} \varphi_A(|A \cap X|)$ (φ_A : concave)
 \mathcal{T} : laminar ($A, B \in \mathcal{T} \Rightarrow A \cap B = \emptyset$ or $A \subseteq B$ or $A \supseteq B$)
- max-value $\mu(X) = \max\{a_i \mid i \in X\}$
- matroid rank (Fujishige 05)
 $\mu(X) = \max\{|I| \mid I : \text{independent}, I \subseteq X\}$
- weighted matroid rank ($w \geq 0$) (Shioura 09)
 $\mu(X) = \max\{w(I) \mid I : \text{independent}, I \subseteq X\}$

Dual Character of Matroid Rank Func

$$\rho(X) = \max\{|I| \mid I : \text{independent}, I \subseteq X\}$$

is **L[♠]-convex** and **M[♠]-concave**

Self-Conjugacy: $\rho(X) = |X| - \rho^\bullet(\chi_X)$

Edmonds' matroid union formula:

$$\max_X \{\rho_1(X) + \rho_2(V \setminus X)\} = \min_Y \{\rho_1(Y) + \rho_2(Y) + |V \setminus Y|\}$$

submod maximization

(M[♠]-concave + M[♠]-concave)

submod minimization

(L[♠]-convex + L[♠]-convex)

Algorithms for Submodular Set Func

convex extension

concave extension

computable as
Lovász extension
(Lovász 03)

subclass: M^{\natural} -concave
(valuated matroid)

1. **Greedy** algorithm for max of **an** M^{\natural} -concave fn
(Dress-Wenzel 90)
2. **Matroid intersection**-type algorithm for max of a
sum of **two** M^{\natural} -concave fns (Murota 96)
3. **Pipage rounding** algorithm for approx max of a
sum of **several** nondecr. M^{\natural} -concave fns (Shioura 09)

M[♯]-concave Maximization Algorithm

$\mu(X)$: M[♯]-concave set function ($\mu(\emptyset) > -\infty$)

Greedy algorithm

S0: Put $X := \emptyset$

S1: Find $j \in V \setminus X$ that maximizes $\mu(X \cup \{j\})$

S2: If $\mu(X) \geq \mu(X \cup \{j\})$, stop (X :maximizer of μ)

S3: Set $X := X \cup \{j\}$ and go to S1

- (Variant of) Dress–Wenzel's for valuated matroid
- Kruskal's algorithm for min spanning tree

Submodular Max. under Matroid Constraint

Maximize $f(S)$ s.t. $|S| = k$

Maximize $f(S)$ s.t. $|S| \leq k$

Maximize $f(S)$ s.t. S : base in a matroid

Maximize $f(S)$ s.t. S : independent in a matroid

Submodular function maximization under a matroid constraint is an NP-hard problem in general

BUT

If f is M^{\natural} -concave, this is an M^{\natural} -concave intersection problem and, hence poly-time solvable

Submodular Welfare Maximization

Maximize $f_1(S_1) + f_2(S_2) + \dots + f_k(S_k)$ s.t.

(S_1, S_2, \dots, S_k) : partition of S (f_i : submodular)

Submodular welfare maximization is an NP-hard problem in general

BUT

If f_i 's are M^\natural -concave, this reduces to an M^\natural -concave intersection (or, M^\natural -concave convolution) and, hence poly-time solvable.

This is the case if $f_i(X) = \varphi_i(|X|)$ with concave φ_i

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